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# Improving performance of contour integral-based nonlinear eigensolvers with infinite GMRES

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Yuqi Liu

School of Mathematical Sciences, Fudan University



Joint work with Jose E. Roman and Meiyue Shao

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# Preliminaries

- Nonlinear eigenvalue problem (NEP):

$$T(\lambda)v = 0, \quad v \in \mathbb{C}^n \setminus \{0\}, \quad \lambda \in \Omega.$$

- $\Omega$ : a connected region with a smooth boundary.
  - $T(\xi)$ : a holomorphic  $\xi$ -dependent matrix in  $\Omega$ .
  - $\lambda$ : eigenvalue.
  - $v$ : (right) eigenvector.
- Our goal:
    - To find all eigenvalues  $\lambda_1, \dots, \lambda_k$  lying in  $\Omega$ ,  
as well as their corresponding (right) eigenvectors  $v_1, \dots, v_k$ .

- Generalized eigenvalue problems:  $(A - \lambda B)v = 0$ . ( $T(\xi) = A - \xi B$ )
- Polynomial eigenvalue problems:  $(A_0 + \lambda A_1 + \lambda^2 A_2 + \cdots + \lambda^p A_p)v = 0$ .
- Quasi normal mode (QNM) analysis for nanophotonic devices:

$$\sum_{j=1}^p r_j(\lambda) L_j v = 0, \quad r_j(\cdot): \mathbb{C} \rightarrow \mathbb{C}, \quad L_j \in \mathbb{C}^{n \times n}, \quad j = 1, \dots, p.$$

–  $r_j(\cdot)$ : rational functions.

- Acoustic analysis of the Helmholtz wave equation:

$$(A + \lambda B + \lambda^2 C)v = e^{i\lambda\tau} S v, \quad \tau \in \mathbb{R}, \quad A, B, C, S \in \mathbb{C}^{n \times n}.$$

- Approximate the moments

$$\mathcal{M}_j = \frac{1}{2\pi i} \int_{\partial\Omega} \xi^j T(\xi)^{-1} Z d\xi,$$

by quadrature rules (for example, the mid point rule)

$$\mathcal{M}_{j,N} = \frac{1}{iN} \sum_{j=0}^{N-1} \varphi^j(\theta_j) \varphi'(\theta_j) T(\varphi(\theta_j))^{-1} Z.$$

And solve some small-size problems with these moments to approximate the eigenpairs.

- The nonlinear Sakurai-Sugiura algorithm.
  - The nonlinear FEAST algorithm.
  - **Beyn's method.**
- Drawback: A lot of large, sparse linear systems have to be solved!

- If  $T(\xi) = I - \xi A$ , we can solve all  $T(\xi_j)^{-1}z$ ,  $j = 0, 1, \dots, N-1$ , with a single Arnoldi process.
  - The Krylov subspaces of  $I - \xi_j A$  are the same.

$$\mathcal{K}_m(I - \xi_j A, z) = \mathcal{K}_m(A, z), \quad j = 0, \dots, N-1.$$

- If we have

$$AU_m = U_{m+1} \underline{H}_m$$

by the Arnoldi process, we also have

$$(I - \xi_j A)U_m = U_{m+1}(\underline{I}_m - \xi_j \underline{H}_m).$$

- The approximate solution can be obtained by solving  $N$  small least squares problems

$$U_m \cdot \arg \min_y \|(\underline{I}_m - \xi_j \underline{H}_m)y - \|z\|e_1\|_2.$$

- To solve  $T(\xi_0)^{-1}z$  for some  $\xi_0$  close to 0, we can transform it into

$$T(\xi_0)^{-1}z \xrightarrow{(1)} \left( \sum_{j=0}^p \frac{\xi_0^j}{j!} T^{(j)}(0) \right)^{-1} z \xrightarrow{(2)} (\mathcal{L}_0 - \xi_0 \mathcal{L}_1)^{-1} \hat{z} = \mathcal{L}_0^{-1} (I - \xi_0 \mathcal{L}_1 \mathcal{L}_0^{-1})^{-1} \hat{z},$$

where

$$\mathcal{L}_0 = \begin{bmatrix} T(0) & \frac{T^{(1)}(0)}{1!} & \frac{T^{(2)}(0)}{2!} & \cdots & \frac{T^{(p)}(0)}{p!} \\ & I & & & \\ & & I & & \\ & & & \ddots & \\ & & & & I \end{bmatrix}, \quad \mathcal{L}_1 = \begin{bmatrix} 0 & & & & \\ I & 0 & & & \\ & I & \cdots & & \\ & & \cdots & 0 & \\ & & & I & 0 \end{bmatrix}, \quad \hat{z} = \begin{bmatrix} z \\ 0 \\ \vdots \\ 0 \end{bmatrix}.$$

(1). Taylor expansion at  $\xi = 0$ .

(2). Companion linearization.

- We can solve **several** different  $\xi_0$ 's with **a single Arnoldi process** on  $\mathcal{L}_1 \mathcal{L}_0^{-1}$ .
- Advantage: **only one factorization** on  $T(0)$  is needed.



**Input:** Maximum iteration  $m$ , the parameter-dependent matrix  $T(\xi): \mathbb{C} \rightarrow \mathbb{C}^{n \times n}$ , the right-hand side  $z \in \mathbb{C}^n$  and the points to be solved  $\xi_j$  for  $j = 0, \dots, N-1$

**Output:** Approximations  $x_{0,j} \approx T(\xi_j)^{-1}z$  for  $j = 0, \dots, N-1$

1: Linearize  $T$  to

$$\mathcal{L}_0 = \begin{bmatrix} T(0) & \frac{T^{(1)}(0)}{1!} & \dots & \frac{T^{(p)}(0)}{p!} \\ & I & & \\ & & \dots & \\ & & & I \end{bmatrix}, \quad \mathcal{L}_1 = \begin{bmatrix} 0 & & & & \\ I & 0 & & & \\ & \dots & \dots & & \\ & & & I & 0 \end{bmatrix}, \quad p > m$$

2: Perform  $m$  iterations of Arnoldi process on  $(\mathcal{L}_1 \mathcal{L}_0^{-1}, \hat{z})$  to obtain  $\mathcal{L}_1 \mathcal{L}_0^{-1} \mathcal{U}_m = \mathcal{U}_{m+1} \underline{H}_m$

3: Set  $y_j \leftarrow \arg \min_y \|(I_m - \xi_j \underline{H}_m)y - \|z\|e_1\|_2$  for  $j = 0, \dots, N-1$

4: Set  $x_j \leftarrow \mathcal{L}_0^{-1} \mathcal{U}_m y_j$  for  $j = 0, \dots, N-1$

5: Set  $x_{0,j} \leftarrow x_j(1:n)$  for  $j = 0, \dots, N-1$

- Remark: when taking  $p > m$ , we can assume there is **no truncate error** introduced by Taylor expansion. In other words, the parameter-dependent matrix  $T(\xi)$  is expanded **infinitely**.

# Accelerate Beyn's method with infinite GMRES

**Input:** The parameter-dependent matrix  $T(\xi): \mathbb{C} \rightarrow \mathbb{C}^{n \times n}$ , the initial guess

$Z = [z_1, \dots, z_k] \in \mathbb{C}^{n \times k}$ , the contour  $\varphi$  and quadrature nodes  $\theta_j$ 's for  $j = 0, \dots, N-1$

**Output:** Approximate eigenvalues  $\Lambda$  and eigenvectors  $V$

1: **for**  $s = 1, \dots, k$  **do**

2: Use infGMRES to solve  $T(\varphi(\theta_j))^{-1} z_s$  for  $j = 0, \dots, N-1$  simultaneously

3: **end for**

4: Set  $\mathcal{M}_{0,N} \leftarrow \frac{1}{iN} \sum_{j=0}^{N-1} \varphi'(\theta_j) T(\varphi(\theta_j))^{-1} Z$

5: Set  $\mathcal{M}_{1,N} \leftarrow \frac{1}{iN} \sum_{j=0}^{N-1} \varphi(\theta_j) \varphi'(\theta_j) T(\varphi(\theta_j))^{-1} Z$

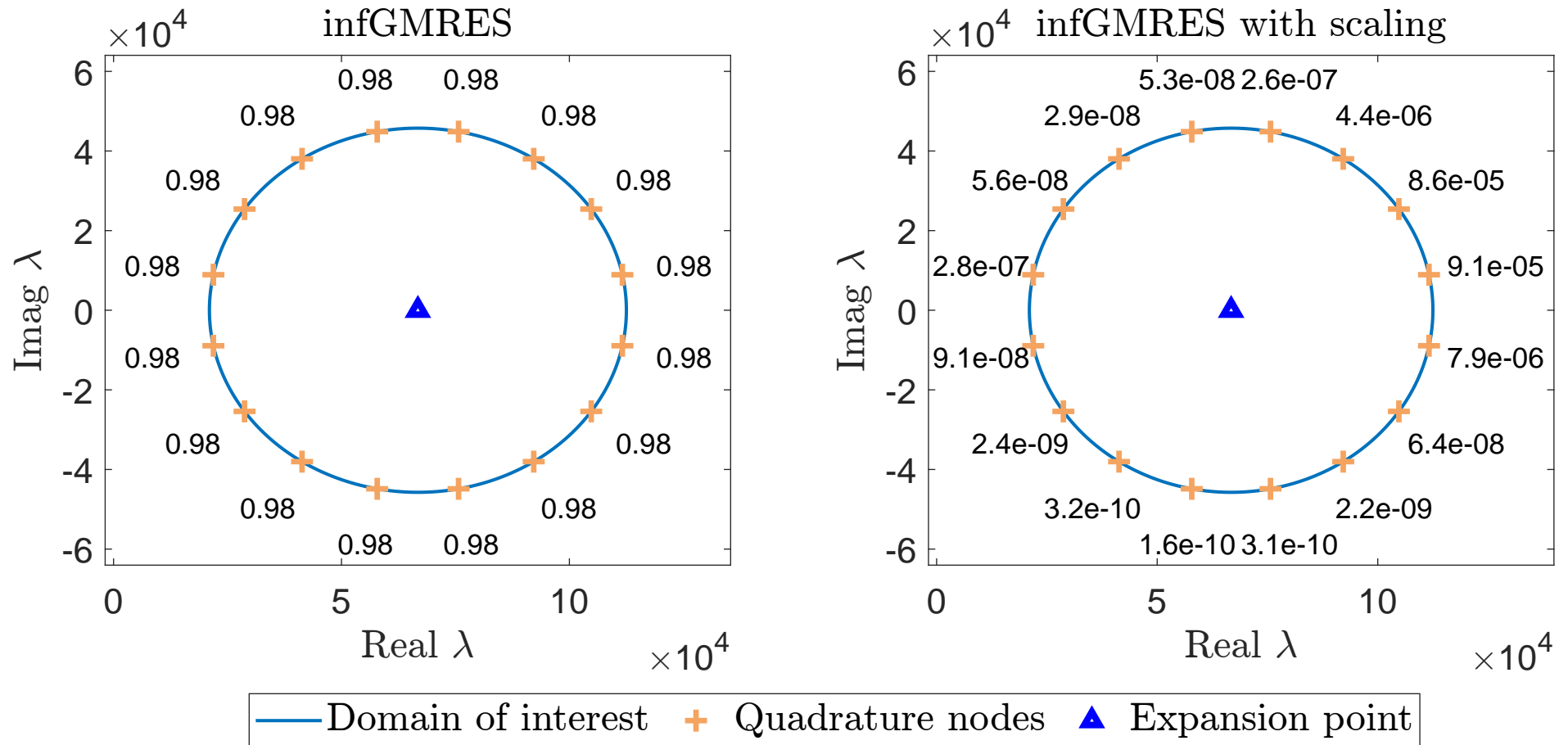
6: Singular value decomposition  $\mathcal{M}_{0,N} = V_0 \Sigma_0 W_0^*$

7: Set  $\check{\mathcal{M}}_{1,N} \leftarrow V_0^* \mathcal{M}_{1,N} W_0 \Sigma_0^{-1}$

8: Eigenvalue decomposition  $\check{\mathcal{M}}_{1,N} = S \Lambda S^{-1}$

9: Set  $V \leftarrow V_0 S$

- The simple implementation is not efficient enough!
  - High accuracy can not be reached with a modest GMRES iteration  $m$ .
  - Storing Arnoldi subspace  $\mathcal{U}_m$  requires  $\mathcal{O}(m^2 n)$  memory.



- The Taylor expansion is employed on the center of a circular contour.
- The relative residuals of 16 linear parameterized systems of the gun problem are shown.
- The accuracy is obviously higher when scaling  $T(a\xi + b)$  is used (right).

- On the scaling technique:
  - Reveal the relationship between NEP and their scaled problems.  
(It does not make sense that a scaling will increase the accuracy of the Taylor expansion.)
  - Provide a novel weighting strategy to accelerate the convergence of infGMRES.
- On the memory usage:
  - Adopt the technique of TOAR to represent the Arnoldi basis compactly (Omitted).

A weighting technique

- Solving  $T(\xi)^{-1}z$  at  $\xi = \xi_0$  is equivalent to solving  $\tilde{T}(\xi)^{-1}z$  at  $\xi = \xi_0/\rho$ , where  $\tilde{T}(\xi) = T(\rho\xi)$ .
- The linearization becomes  $(\tilde{\mathcal{L}}_0 - \frac{\xi_0}{\rho}\tilde{\mathcal{L}}_1)^{-1}\hat{z}$ , where

$$\tilde{\mathcal{L}}_0 = \begin{bmatrix} T(0) & \rho \frac{T^{(1)}(0)}{1!} & \rho^2 \frac{T^{(2)}(0)}{2!} & \dots & \rho^p \frac{T^{(p)}(0)}{p!} \\ & I & & & \\ & & I & & \\ & & & \ddots & \\ & & & & I \end{bmatrix}, \quad \tilde{\mathcal{L}}_1 = \mathcal{L}_1.$$

- We notice that

$$\tilde{\mathcal{L}}_0 = D_\rho^{-1} \mathcal{L}_0 D_\rho, \quad \frac{1}{\rho} \tilde{\mathcal{L}}_1 = D_\rho^{-1} \mathcal{L}_1 D_\rho, \quad \tilde{\mathcal{L}}_0 - \frac{\xi_0}{\rho} \tilde{\mathcal{L}}_1 = D_\rho^{-1} (\mathcal{L}_0 - \xi_0 \mathcal{L}_1) D_\rho,$$

where

$$D_\rho = \begin{bmatrix} I & & & \\ & \rho I & & \\ & & \ddots & \\ & & & \rho^p I \end{bmatrix}.$$

- Actually, more degrees of freedom can be introduced.

**Lemma 1** Suppose  $d_j \in \mathbb{C} \setminus \{0\}$  and  $T_j \in \mathbb{C}^{n \times n}$  for  $j = 0, \dots, p$ , and

$$D = \begin{bmatrix} d_0 I & & & & \\ & d_1 I & & & \\ & & \dots & & \\ & & & \dots & \\ & & & & d_p I \end{bmatrix}, \quad \mathcal{L}_0 = \begin{bmatrix} T_0 & T_1 & T_2 & \dots & T_p \\ & I & & & \\ & & I & & \\ & & & \dots & \\ & & & & I \end{bmatrix}, \quad \mathcal{L}_1 = \begin{bmatrix} 0 & & & & \\ I & 0 & & & \\ & \dots & \dots & & \\ & & & I & 0 \end{bmatrix}.$$

Then, for any scalar  $\xi \in \mathbb{C}$  and vector  $z \in \mathbb{C}^n$ ,

$$(D^{-1}(\mathcal{L}_0 - \xi \mathcal{L}_1)D)^{-1} \hat{z} = \begin{bmatrix} (\sum_{j=0}^p \xi^j T_j)^{-1} z \\ * \\ \vdots \\ * \end{bmatrix}.$$

- Therefore, the scaling in infGMRES can be regarded as a balanced companion linearization  $(D^{-1}\mathcal{L}_0D, D^{-1}\mathcal{L}_1D)$ .



- If the Krylov subspace  $\mathcal{K}_m(I - \xi_0 \mathcal{L}_1 \mathcal{L}_0^{-1}, \hat{z})$  is spanned by  $\mathcal{U}_m$ , by induction, we can prove that  $\mathcal{K}_m(D^{-1}(I - \xi_0 \mathcal{L}_1 \mathcal{L}_0^{-1})D, \hat{z})$  is spanned by  $D^{-1}\mathcal{U}_m$ .
- At the  $m$ th iteration, infGMRES will give the solution by

$$\begin{aligned} y_* &= \arg \min_y \left\| (D^{-1}(I - \xi_0 \mathcal{L}_1 \mathcal{L}_0^{-1})D)(D^{-1}\mathcal{U}_m)y - \hat{z} \right\|_2 \\ &= \arg \min_y \left\| D^{-1}((I - \xi_0 \mathcal{L}_1 \mathcal{L}_0^{-1})\mathcal{U}_m y - \hat{z}) \right\|_2. \end{aligned}$$

- Idea: we can choose  $D$  appropriately to guide GMRES to pick a more accurate solution from the search space!
  - Our goal is to minimize  $\|r_N\|_2 = \|T(\xi_0)x_0 - z\|_2$ , or  $\|r_P\|_2 = \left\| \sum_{j=0}^p \xi_0^j T_j x_0 - z \right\|_2$ , where  $x_0 = (D^{-1}\mathcal{L}_0^{-1}\mathcal{U}_m y_*)(1:n)$ .

**Lemma 2** Suppose vector  $y_* \in \mathbb{C}^m$  is the solution for some weighting matrix  $D$ . Then, the polynomial-wise residual  $r_P = \sum_{j=0}^p \xi_0^j T_j x_0 - z$  can be represented as

$$r_P = \left[ I, -\sum_{j=1}^p \xi_0^{j-1} T_j, -\sum_{j=2}^p \xi_0^{j-2} T_j, \dots, -T_p \right] \begin{bmatrix} r_0 \\ r_1 \\ \vdots \\ r_p \end{bmatrix},$$

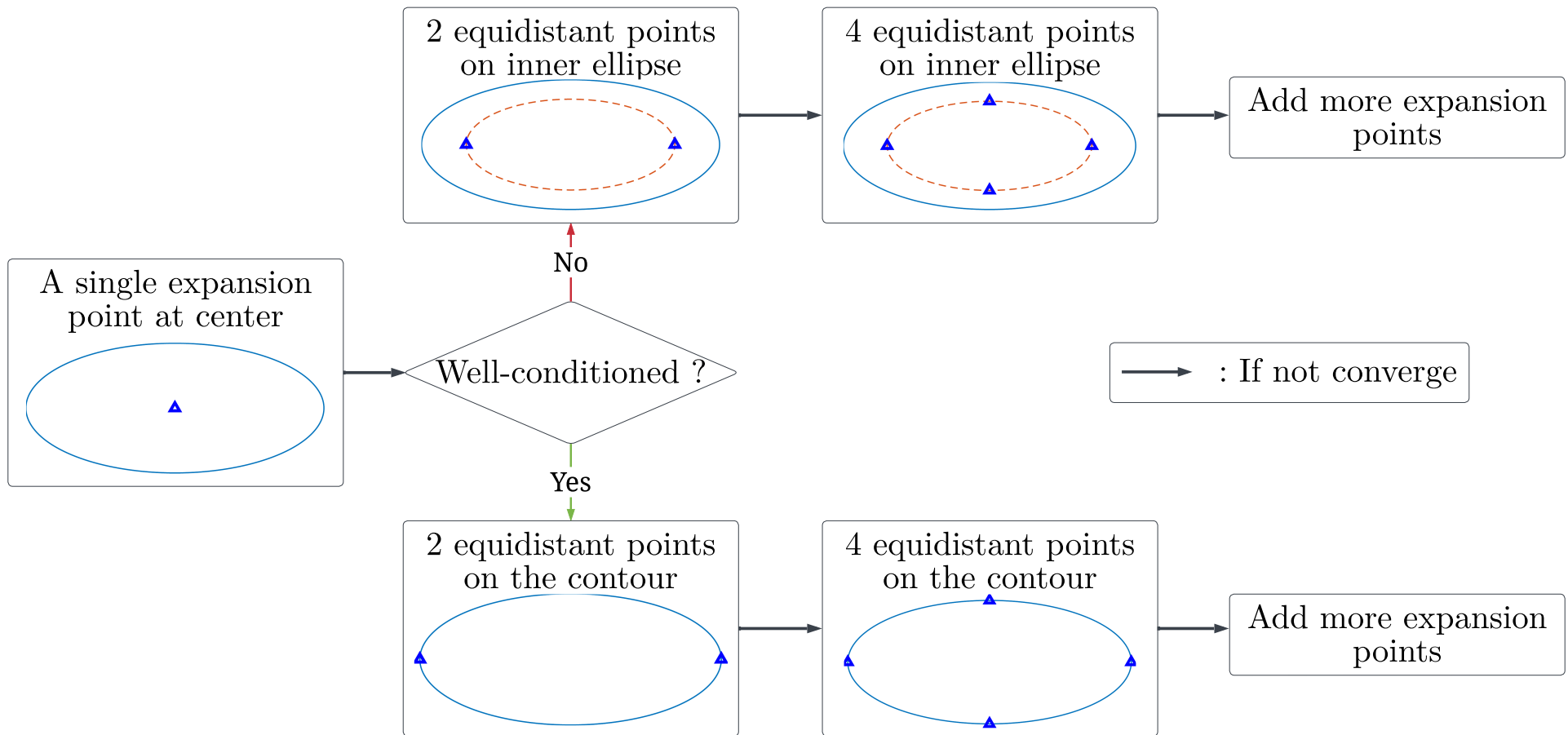
where

$$\begin{bmatrix} r_0 \\ r_1 \\ \vdots \\ r_p \end{bmatrix} = (I - \xi_0 \mathcal{L}_1 \mathcal{L}_0^{-1}) \mathcal{U}_m y_* - \hat{z}.$$

- $r_j$  can be adjusted by  $D$ .
- We can decrease the residual by balancing the sum

$$\|r_P\| \leq \|r_0\|_2 + \left\| \sum_{j=1}^p \xi_0^{j-1} T_j \right\|_2 \|r_1\|_2 + \dots + \|T_p\|_2 \|r_p\|_2.$$

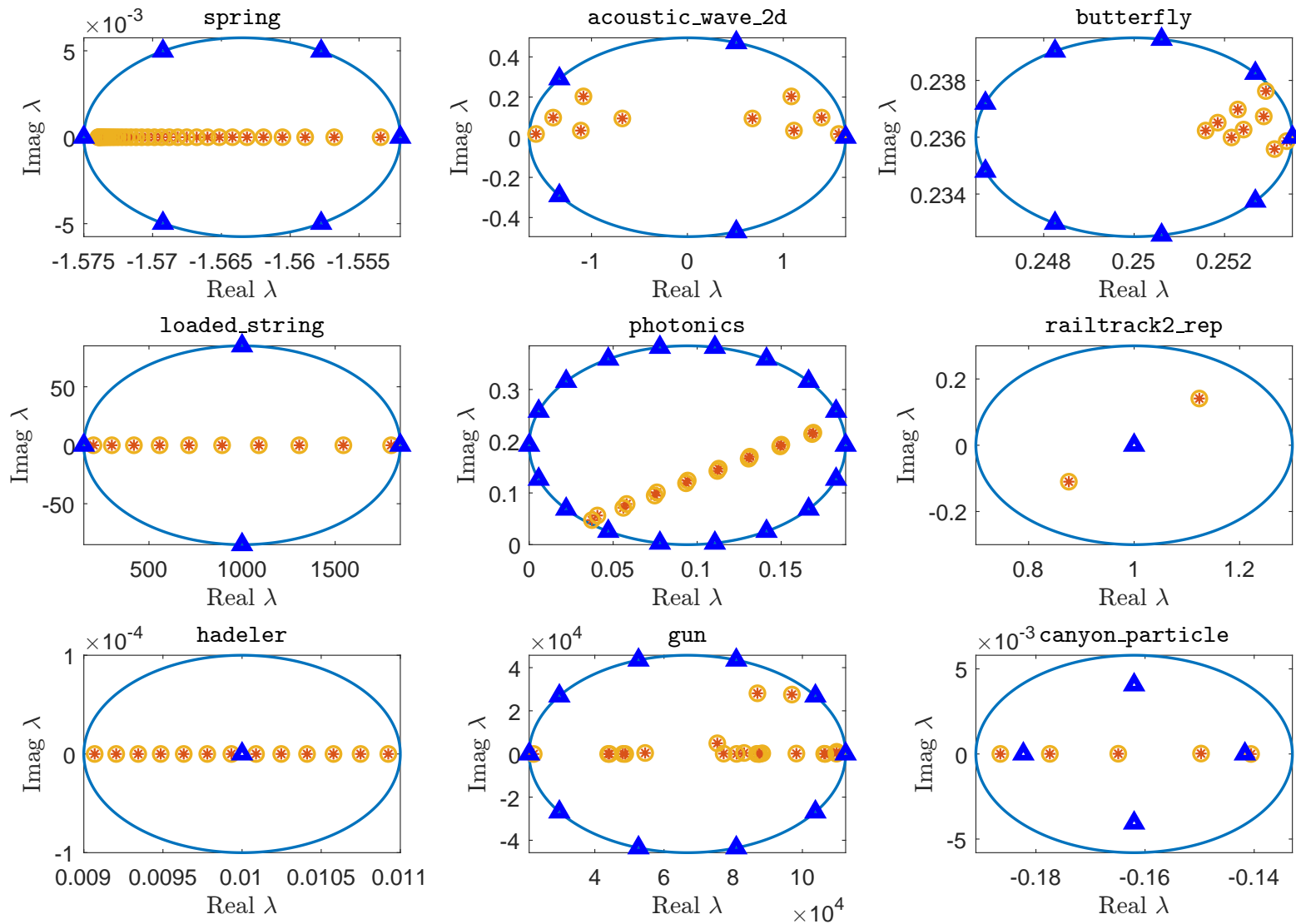
How to choose expansion points



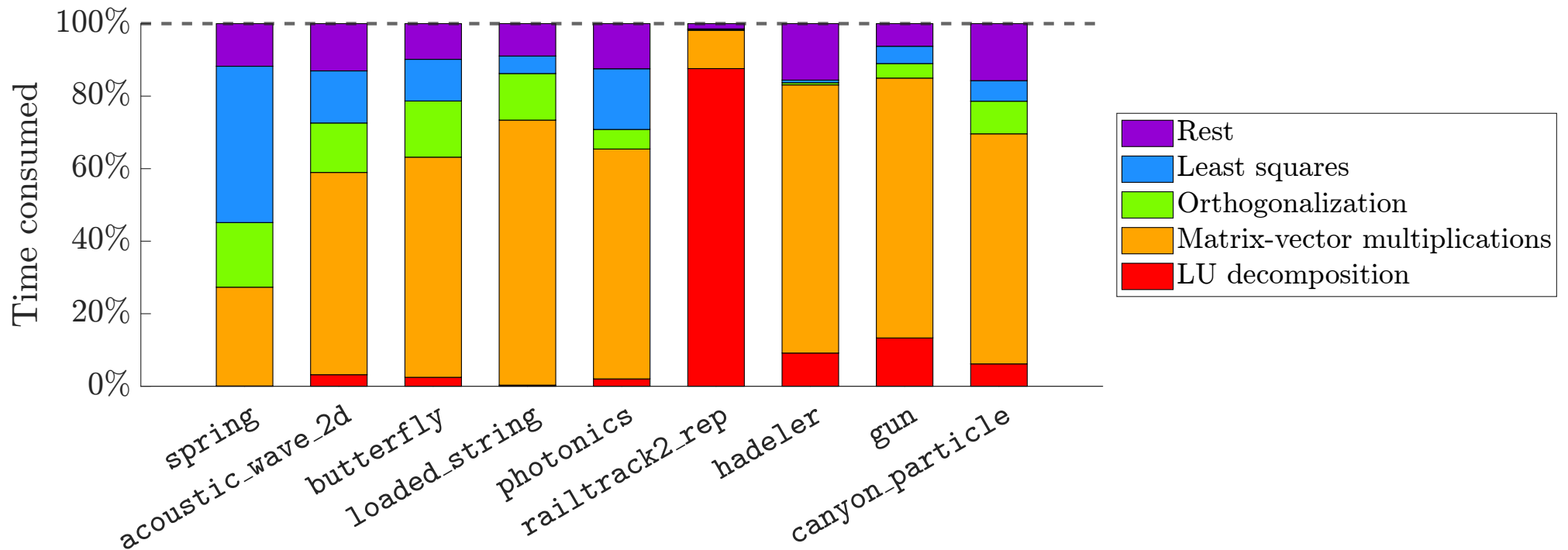
# Numerical Experiments

- Information of test problems.
  - $n$ : the size of the problem.
  - $k$ : the number of the eigenvalues to be computed.
  - $N$ : the number of quadrature nodes.
  - $n_{ep}$ : the number of expansion points.
  - $t_s, t_{iG}$ : the times consumed by Beyn's method with MATLAB backslash and with infGMRES.

Problem	Type	$n$	$k$	$N$	$n_{ep}$	$t_s$ (s)	$t_{iG}$ (s)	$(t_s - t_{iG})/t_s$
spring	QEP	3000	32	1024	6	3.472	15.72	−353%
acoustic_wave_2d	QEP	9900	10	512	5	28.14	9.475	66%
butterfly	PEP	5000	9	512	9	11.54	8.453	27%
loaded_string	REP	20000	10	128	4	1.07	9.805	−816%
photonics	REP	20363	16	3060	18	600	209.7	65%
railtrack2_rep	REP	35955	2	128	1	533.4	71.95	87%
hadeler	NEP	5000	13	32	1	40.2	46.6	−16%
gun	NEP	9956	21	1024	10	501.2	148.9	70%
canyon_particle	NEP	16281	5	256	4	40.94	7.652	81%



— Domain of interest  
 \* Approximate eigenvalues  
 ○ True eigenvalues  
 ▲ Expansion points



- Most of time is consumed by the Arnoldi process (matrix–vector multiplications and orthogonalization).
- When the size of the matrices grow higher, the time consumed by matrix factorization increases dramatically.



- Summary

- infGMRES can be employed in contour integral-based nonlinear eigensolvers to reduce the cost of solving linear systems.
- We propose the convergence-accelerating weighting strategy, the memory-friendly TOAR-like trick, and the selection strategies of the parameters, making infGMRES robust and efficient in practice.
- Our algorithm can achieve a speedup up to 7x on the test examples.

- Ongoing work

- Develop a machine learning-based adaptive strategy for selecting expansion points automatically.
- Implement a block variant of infGMRES.

Thank you for your attention!

# Appendix

# Compact representation of the Arnoldi basis

- The Arnoldi process on  $\mathcal{L}_1\mathcal{L}_0^{-1}$  can be rearranged into

$$\mathcal{L}_1\mathcal{L}_0^{-1}\mathcal{U}_m = \mathcal{L}_1\mathcal{L}_0^{-1} \begin{bmatrix} Q_m \check{U}_{m,0} \\ \vdots \\ Q_m \check{U}_{m,p} \end{bmatrix} = \begin{bmatrix} Q_m \check{U}_{m,0} & [Q_m, q_{m+1}] \check{u}_{m+1,0} \\ \vdots & \vdots \\ Q_m \check{U}_{m,p} & [Q_m, q_{m+1}] \check{u}_{m+1,p} \end{bmatrix} \underline{H}_m = \mathcal{U}_{m+1} \underline{H}_m,$$

where

$$Q_m \in \mathbb{C}^{n \times m}, \quad Q_m^* Q_m = I, \quad \check{U}_{m,j} \in \mathbb{C}^{m \times m}, \quad j = 0, \dots, p.$$

- Then, we can set

$$Q_{m+1} \leftarrow [Q_m, q_{m+1}], \quad \check{U}_{m+1,j} \leftarrow \left[ \begin{array}{c|c} \check{U}_{m,j} & \\ \hline 0 & \check{u}_{m+1,j-1} \end{array} \right].$$

- To store  $\mathcal{U}_m$  in this way, we need to store  $Q_m$  by  $\mathcal{O}(mn)$  memory and  $\check{U}_{m,j}$  for  $j = 0, \dots, m-1$  by  $\mathcal{O}(m^3)$  memory.
- When  $n$  is very large,  $\mathcal{O}(mn + m^3) \ll \mathcal{O}(m^2n)$ .

# Comparison on different weighting strategies

- No weight:  $D = I$ .
- Scalar weight from [Betcke 2009]:

$$\rho = \left( \frac{\|T(0)\|_2}{\|T^{(p)}(0)/p!\|_2} \right)^{1/p}, \quad D = \begin{bmatrix} I & & & \\ & \rho I & & \\ & & \dots & \\ & & & \rho^p I \end{bmatrix}.$$

- Our weight:

$$d_0 = 1, \quad d_s = \frac{\gamma}{\|\sum_{j=s}^p \xi_0^{j-s} T_j\|_2}, \quad \gamma = \frac{\|\sum_{j=2}^p \xi_0^{j-2} T_j\|_2^2}{\|\sum_{j=3}^p \xi_0^{j-3} T_j\|_2}, \quad s = 1, \dots, p,$$

$$D = \begin{bmatrix} d_0 I & & & \\ & d_1 I & & \\ & & \dots & \\ & & & d_p I \end{bmatrix}.$$

